

Cerebro: Presence Protocol for Large Mesh Networks

Polychronis Ypodimatopoulos
Viral Communications group
MIT Media Laboratory

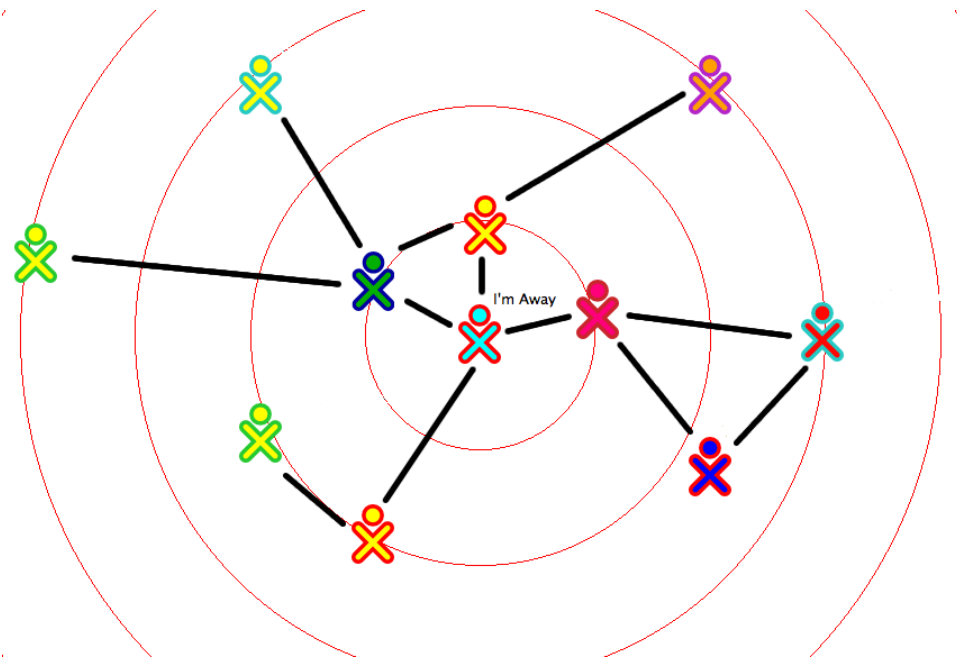
April 2008

Project goal

- Organize the presence, profile and social interaction of humans and objects in the same physical area and make this information accessible and useful

Presence Protocol Overview

- A node receives presence beacons from its neighbors, eliminates duplicate information, creates a new beacon that combines information the node already has with information it received and broadcasts a new beacon to its own neighbors.



Beacon Frame contents

Node ID

Witness

Distance

Serial No.

ID of the node this entry is about
(MAC address, 6 bytes)

ID of the node from which this entry was received
(MAC address, 6 bytes)

Distance from the current node to <Node ID>
(similar to ETX metric)
(1 byte)

Unique ID for each entry published by node <Node ID>
(1 byte)

Beacon Frame contents

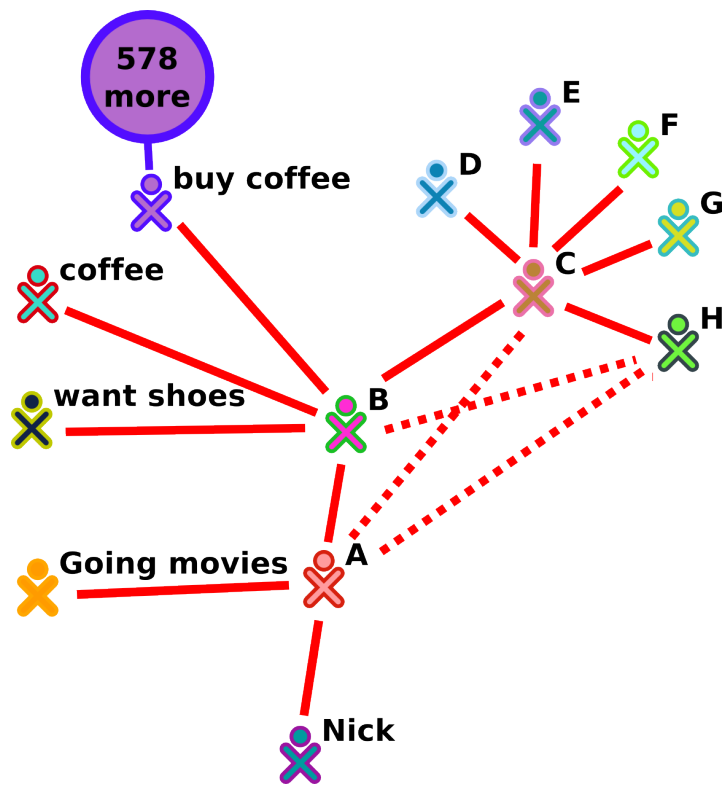
Presence table at **A**:



| Node ID | Witness | Distance |
|---------|---------|----------|
| B | A | 1 |
| C | B | 1 |
| C | A | 2 |
| D | B | 3 |
| E | B | 3 |
| F | B | 3 |
| G | B | 3 |
| H | B | 3 |
| H | A | 5 |

Beacon Frame contents

Presence table at **B**:



| Node ID | Witness | Distance |
|---------|---------|----------|
| A | B | 1 |
| C | B | 1 |
| D | C | 2 |
| E | C | 2 |
| F | C | 2 |
| G | C | 2 |
| H | C | 2 |
| H | B | 2.5 |
| C | A | 3 |

Presence Protocol

- 1) Wait for a period T for beacons from neighboring nodes
- 2) For every beacon received:
 - ✓ For every entry in beacon:
 - if it is about the current node, or the current node is the witness, or the the serial number is not newer than the existing one, discard it!
 - if node/witness pair exists in presence table update number of arrivals and next arrival estimate both for node and witness
 - else add new entry in the table
- 3) Eliminate stale entries in presence table (ie. entries where the next arrival estimate has lapsed)
- 4) Create a new beacon using the minimum distances to each node in presence table
- 5) Broadcast beacon to neighbors

Arrival estimates

- 1) Count arrivals over time period T .
- 2) Formulate a Poisson arrival rate (assuming arrivals are independent events)
- 3) Estimate time of next arrival with 90% accuracy

Example:

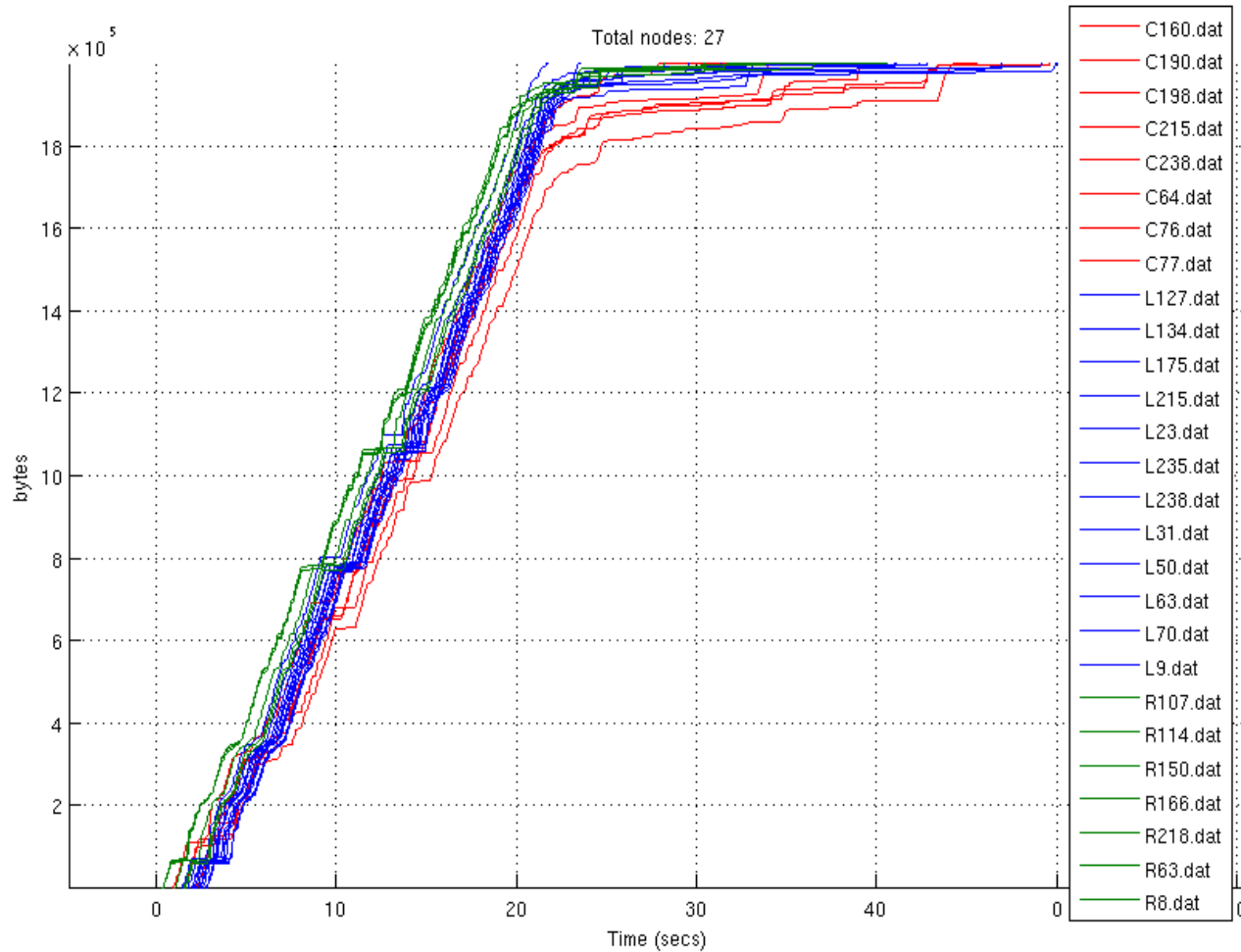
For $T=1\text{sec}$, accuracy=90%, next arrival in 2.3sec

For $T=1\text{sec}$, accuracy=100%, next arrival in infinity (!)

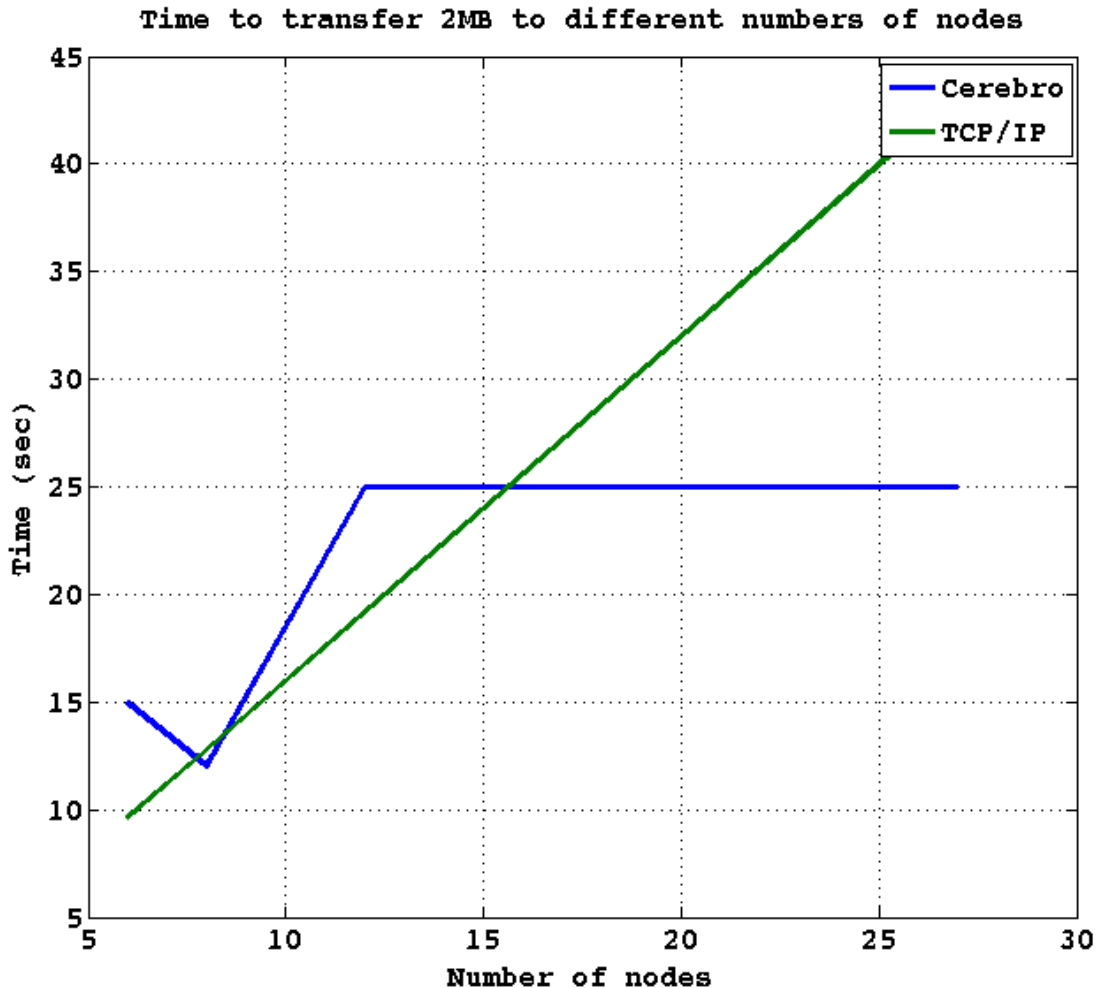
Features

- Achieving scalability “on a diet”: Connected 100 nodes in mesh network using a single frame per node, per 10 seconds (15kb/sec in the worst case)
- Adjustable beacon rate based on node mobility
$$\text{Overhead} = N \cdot B / T = 15 \cdot N^2 / T = 15 \cdot N / t, \quad T = N / t, \quad t = 10\text{Hz}$$
- Different beacon rates in different areas of the mesh network, according to node concentration
- Portability: Cerebro runs on x86, OLPC XO, Nokia N800 and ARM-based embedded computers (python)
- Routing protocol for communication with any other node in the network

Numbers 1/2



Numbers 2/2



Cerebro: 1Mbps
(broadcast)
simultaneous
transfer.

TCP/IP: sequential,
10Mbps (actual)
transfer is assumed
to target nodes.

Interfacing with Cerebro

- Based on DBus
- Methods currently with DBus interface:
 - `register`: register activities/apps
 - `push_data(destinations, data, port)`
 - `onNodeArrival(node_array)`
- Methods without DBus interface (yet):
 - `set_status_info, get_status_info`
 - `request_data(dest, port, request="")`
 - `request_multicast_data(destinations, port, request="")`
 - `push_multicast_data(destinations, port, data)`
 - `handle_new_req(src, port, req_payload)`
 - `handle_new_data(data, fhash, src, port)`
 - `onNodeLeave`

Demo (or Die!)

- Chat
- `/tree`: shows network tree rooted at your XO
- `/sendfile`
- `/getstats`
- `/savestats`